

ETH

Eidgenössische Technische Hochschule Zürich
Swiss Federal Institute of Technology Zurich

Probabilistic embedded systems design



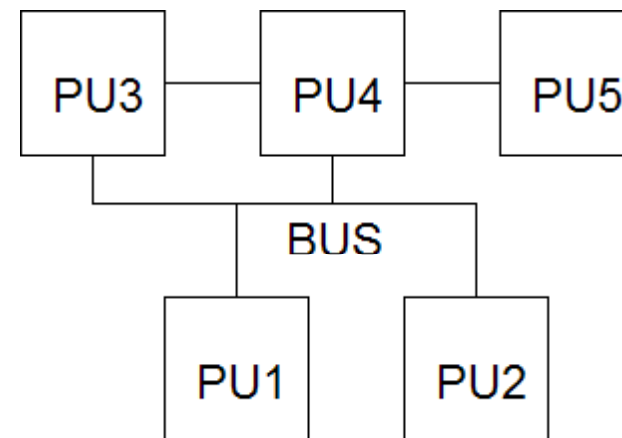


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MPSoC

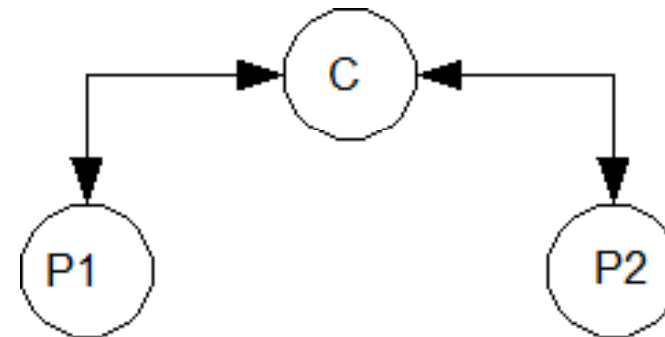
- Multiple processing unit
- Multiple interconnections
- Multiple memory



Hardware Model

$$H = (P, C, L, \varphi, \lambda, \tau, \sigma)$$

P	processing Units
C	communication units
L	edges
φ	power dependent
λ	available comp. resources
τ	bandwidth
σ	delay



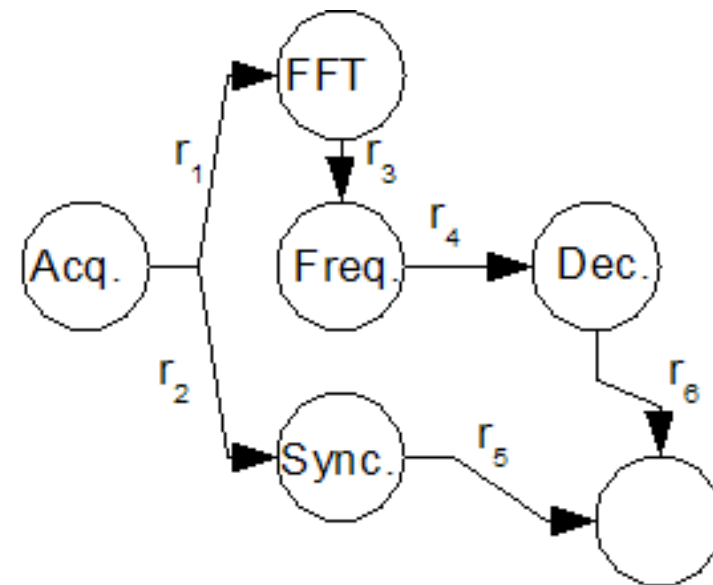
Applications

- Defined by tasks
- Tasks communicate
- Task have an order of execution
- Communication and Computational properties have modes

Application Definition

$$A = (T, E, \gamma, \varepsilon, \theta, \delta, \beta, \omega)$$

T	Tasks
E	Edges
γ	computational compl.
ε	deadline
θ	app. that share task
ω	rate
δ	delay (edge)
β	bandwidth(edge)





Optimization

- Guarantee throughput
- Guarantee delay
- Minimize energy consumption
- Minimize Hardware resources
- Maximize number of applications

Assumptions / Application relation

■ Scenarios

- Set of Applications (\hat{A})
- Group of Applications that share tasks \tilde{A}
- Only one application out of a group \tilde{A} can be active
- Applications in \hat{A} have specific relations
 - $\hat{A}_1 * \hat{A}_2 = 0$
- \hat{A} has a probability of execution
 - Future: probability of execution with notion of time

What's the System

- Tasks interact with the Hardware
 - Tasks are mapped to a specific processing unit
 - Proc. Unit have properties that change tasks behavior
- Robustness ?
- Adaptive ?
- Resilient ?
- Exposure: the set of Appl. defined by Scenarios

Probabilistic approach (1)

- Do not design for worst case scenario
- Define probable use-cases
- Probabilities and Exclusion Rules
 - Reachability of conditions

Probabilistic approach (2)

- Optimize according to scenario

$$\min \left(\sum_{\forall t_i \in T} \sum_{\forall p_j \in P | \text{map}(t_i) = p_j} \left(\frac{\gamma_i \omega_i}{\lambda_j} \text{prob}(\theta_i) \right) \phi_{j0} + \left(1 - \frac{\gamma_i \omega_i}{\lambda_j} \text{prob}(\theta_i) \right) \phi_{j1} \right)$$



Advantages

- less hardware, more applications
- smaller proc. units, less energy consumption
- sufficient for “Soft” real-time systems

Problems

- Analytical proof ?
- Design space explosion
- “Hard” Real-Time systems ?
 - Safety critical devices (Airbus A380)

Related Work

- [1] N. H. Zamora, X. Hu, R. Marculescu, *System-Level Performance/Power Analysis for Platform-Based Design of Multimedia Applications*, ACM Transactions on Design Automation of Electronic Systems, Vol. 12, No. 1, Article 2, January 2007
- [2] S. Hua, G. Qu, S.S. Bhattacharyya, *Energy Reduction Techniques for Multimedia Applications with Tolerance to Deadline Misses*, Proceedings of the Design Automation Conference, pages 131-136, Anaheim CA, June 2003
- [3] J. Hu, R. Marculescu, *Energy- and Performance-Aware Mapping for Regular NoC Architectures*, IEEE Transactions on Computer-Aided Design of Integrated Circuits and Systems, Vol. 24, No. 4, April 2005

A blue-tinted photograph of a large, classical-style building with a prominent dome and arched windows, set against a landscape background.

Questions ?

Thanks for your attention